



## Secure Multiple Multicast in Wireless Network

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**Abstract**-With the emergence of diverse group based services,multiple multicast coexist in single network,user may subscribe to multiple group.However the existing group key management,aiming to secure communication with a single group,are not suitable in multiple group because of inefficient keys,and much large rekeying overhead.In this paper,propose a new GKM for multiple multicast group,called master key encryption based multipl multiple group key management scheme(MKE-MGKM).This MKE\_MGKM exploit asymmetric key that is master key and multiple slave keys,which are generated from proposed master key encryption(MKE)algorithm and is used for efficient distribution of group key.It alleviates rekeying overhead by using asymmetry of master and slave keys. That is even if one of the slave key is updated,the remaining one can be still unchanged by modifying the master key.Through numerical analysis and simulation it is shown that MKE-KGKM can reduce storage overhead of key distribution centre by 75% and storage overhead of user by 85% and 60% of communication overhead,compared to the existing scheme.

**Keywords**-Security, groupkey management, multicast, Chinese remainder theorem, master key encryption

### I. INTRODUCTION

Multicast is an efficient method in transmitting the data from single source to several destination. Especially, in wireless network using a broadcast medium, a single transmission can be received by all nodes within a transmission range, which make it easy to implement the multicast. Therefore, multicast paves efficient way for multiple group communications, by which many group based application, such as charged video on demand or video conferencing, can be commercialized.

The broadcasting medium, however, makes the wireless network to various security attack since anyone can eavesdrop message transmitted in the air.

To implement the multicast that is delivery of message to the member of the group, in wireless network, we need to have an access control mechanism for broadcasted message, which guarantees confidentiality, protects digital contents, and facilitate accurate accounting. Therefore, it is one of the key requirement for successful commercialisation of the multicast service in wireless networks.

The usual way to provide an access control mechanism for secure group communication is by using symmetric key, known as group key, shared only by group members. Messages, encrypted by a member having a group key, can be decrypted by the another group member having the same group key, which can guarantee secure group communication. Although this mechanism, using the shared group key,is an efficient way to guarantee security, it causes difficulties in maintaining efficient key management scheme since the user must be updated before leaving or joining, which is referred as rekeying. To reduce the key management overhead from rekeying, a tree based group key management (GKM) have been studied.

However the existing GKM still faces the limitation of rekeying as the no of multicast services increases. However, in the foreseeable future, multiple multicast group will exist in a single network due to the emergence of many group based application. In such a situation, service provider suffer from the key management overhead for supporting multiple multicast groups. For example, service provider may provide three different multicast in wireless network such as charged TV streaming, a telematics service and an information service. In this example, there will be three user groups for three services, each of which should be managed according to the membership record of the service provider. Moreover, if the subscription for the channel is either charged for each channel or content, the service provider must manage additional user group, for accurate accounting.



## II. OVERVIEW

The usual way to provide access control mechanism for secure group communication is to employ symmetric key, known as group key, shared only by group members. Messages, encrypted by a member having the group key, can be decrypted by another group key having the same group key, which can guarantee secure group communication.

Although this mechanism, using the shared group key, is an efficient way to guarantee security, it causes difficulties in maintaining key management system since the group key must be updated according to the membership change such as user member leaving or joining referred as rekeying. However, the existing GKM still faces the limitation of group key as the no of multicast service increases.

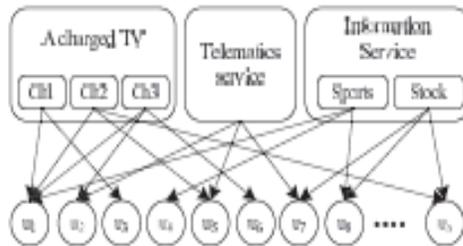
A new group key management called multiple multicast group called master key encryption based multiple group key management (MKE-MGKM) is used in the proposed system.

The MKE-MGKM exploits asymmetric key that is master key and multiple slave key, which are generated from the proposed master key encryption algorithm, and is used for the efficient distribution of group key.

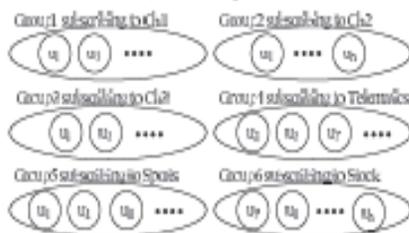
It alleviates the rekeying overhead by using asymmetry of master and slave key, that is even if one of the slave key gets updated, the remaining ones will be still unchanged by modifying only the master key.

The Master key Management Algorithm created and updates the master key and multiple slave key. A message encrypted by master key can be decrypted by several slave key, and vice versa. The key feature of MKE is that one slave key can be revoked by updating master key with simple computation whereas the other key kept valid. This enable us to reduce the rekeying overhead.

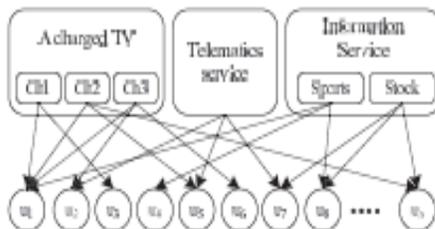
The Rekey for presenting multiple group using the MKE-based key graph having master key and multiple slave key



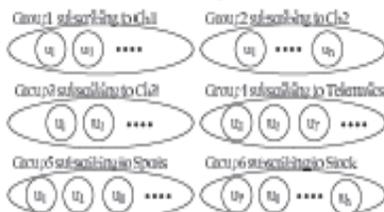
(a) Subscription



(b) Six multicast groups

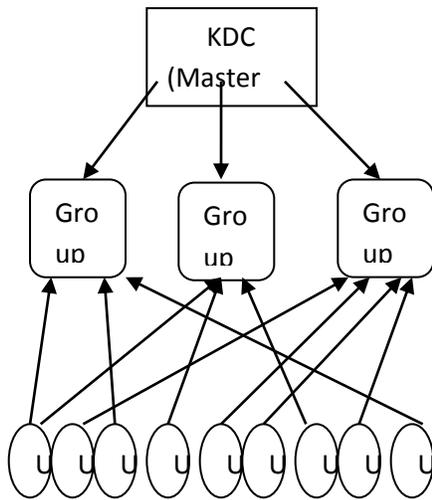


(a) Subscription



(b) Six multicast groups

*Proposed Architecture Diagram*



**III. RELATED WORK**

*Group Key Management Protocol: A Novel Taxonomy*

Group Key Management is the important functional group for multicast services. The invention of internet leads to various services, combining voice, video and text over IP. Although unicast is predominant so far, for media services. Indeed, multicasting is useful for group oriented and in video conferencing, interactive group games, video on demand, TV over Internet, elearning, software updates and database replication. However lack of security cause lot of problem in multicasting. These issue with certain ranges. In public broadcasting, while authentication is a fundamental requirement, confidentiality may not be.

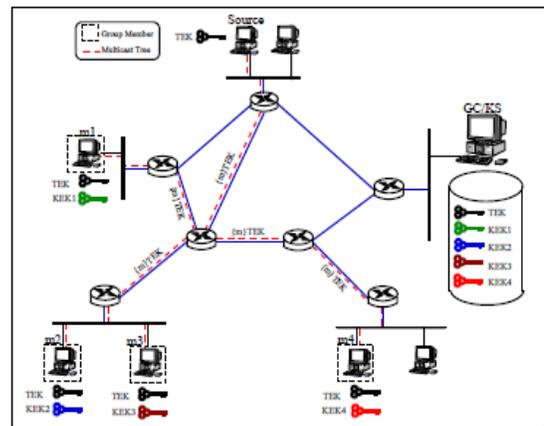
In the contrary case, both confidentiality and authentication are required in video conferencing application. In this paper focus on key store management over wired network:

*Group Key Management.*

The limitation in research have caused lot of confusion in multicast, such as confidentiality, authentication, watermarking and access control

*Group Communication Confidentiality*

The challenging feature for confidentiality. we select a source that sends data to set of receivers in multicast services. The Security is managed by two sessions: Group controller (GC) required for authentication, authorisation and access control, key server required for distribution of keys. depending on key management architecture.



**Simple scenario of Group Management**

To ensure confidentiality of Group Management, the sender should share with all valid members, called traffic encryption key (TEK). To multicast a secret message, it should encrypt with symmetric key algorithm. Upon receiving multicast encrypted message, each valid member can decrypt with TEK and recover the original one. To avoid leaving or rejected member from the group, continue to decrypt the secret message, the KS must distribute TEK and distribute the valid members except the leaving one. This operation is called rekeying. The KS must distribute the Key Encryption Key (KEK) to all member. To rekey, a leave from the group, the KS generates a new TEK: TEK' sends to each member sends it to corresponding KEK. The leaving member doesn't know KEK' and cannot decrypt future multicast messages.

When a new member joins a group, it must be authenticated by the GC. After checking the access of group, they introduce rekey which avoids decrypting of messages using an old member who is not active.



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Therefore KS generates a new TEK: new TEK' encrypts with a old TEK, and it multicast to the group. Hence all old member could recover the new key. so it sends to TEK' to decrypt the message.

The encryption and distribution of key in rekeying is referred as Group Key Management. In this illustrative protocol, rekey induce a new key after leaving of a member(n) where n is the group members. It induce a storage of  $O(n)$  during the multicast services. Since each membership change needs an updation its highly dynamic, without using storage overhead. Proposed architecture in this literature, best storage performance.

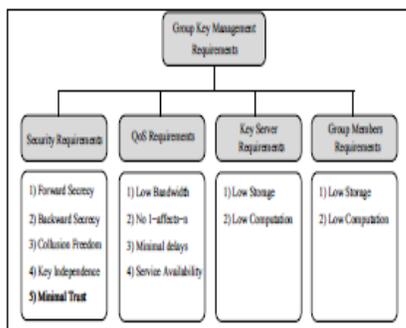
#### Group Key Management Requirement

Efficient Group must manage with the following requirement. These requirement from four point of view.

- Security
- Quality of service
- KS resource and
- Group member resource

It is represented in the following

#### Diagram



Group Key Management Requirement

#### Security Requirement

1. Forward Secrecy who left the group have no access of new group. This ensures that group member leaving cannot decrypt using the old key and enter the future enhanced work. To assure forward secrecy, after leave a new TEK is the new solution.

2. Backward Secrecy ensures that new user must not have access over the old records. It ensures that data must not encrypt before joining a group. To ensure the background secrecy, after join of new TEK is the ultimate solution.
3. Collusion Freedom ensures that any fraudulent user should use the traffic encryption key
4. Key Independence: a protocol is said to be key Independent, if a discloser does not compromise other key.
5. Minimal Trust: The key management must not trust any other entities. Otherwise, the deployment of key would not be easy

#### Quality Of Service Requirement

1. Low Bandwidth Overhead: the rekey must not ensure a group of messages, especially for dynamic groups. Ideally, this should be independent of group size.
2. 1-affects-n: a protocol affects 1-affects-n since one group change will affect the number of group in a message. This affects the new group key member should be updated.

## IV. PROPOSED SYSTEM

In this paper, we propose a new group key management scheme for multiple multicast groups, called the master-key-encryption-based multiple group key management (MKE-MGKM) scheme.

The MKE-MGKM scheme exploits asymmetric keys, i.e., a master key and multiple slave keys, which are generated from the proposed master key encryption (MKE) algorithm, and is used for efficient distribution of the group key.

It alleviates the rekeying overhead by using the asymmetry of the master and slave keys, i.e., even if one of the slave keys is updated, the remaining ones can still be unchanged by modifying only the master key.

#### Modules:

- 1) User Enrollment
- 2) Master Key Management Algorithm
- 3) Rekeying at membership change



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### User Enrollment:

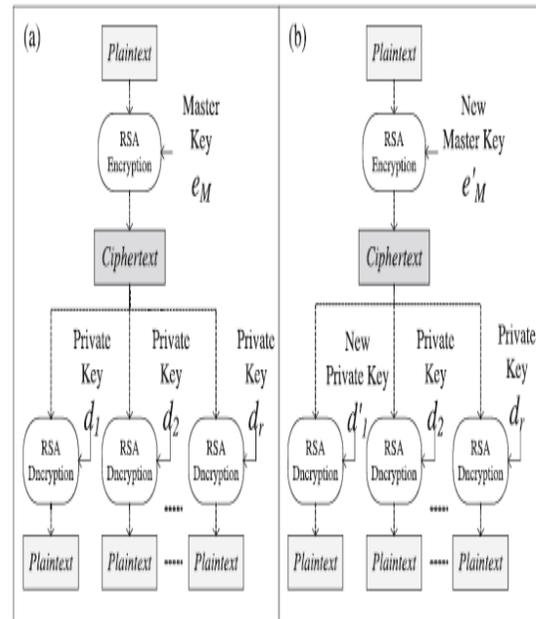
Each User should register their personal details and submit to the server, after submitting server provides the multiple services to access.

### Master Key Management Algorithm:

Multicast Broadcast Services are provided in a network, First of all, the KDC generates a master key, and as many public-private key pairs as the number of Service Group (SGs), through the new proposed master key management algorithm. The number of SGs grows exponentially as the number of Data Group (DGs) increases.

### Rekeying at membership change:

In order to describe the rekeying process of the MKE MGKM scheme, consider the case that a user changes membership. He has been subscribing to particular *MBS1*, but wants to subscribe to *MBS2* instead of *MBS3* while keeping *MBS1*. Therefore, the KDC should switch user's membership from *SG* to particular *SG*, which makes user lose the access privilege to *MBS1* but gain access to *MBS2*. Since the data of *MBS1* should not be visible to that user any more after he un-subscribed from *MBS1*. For the forward secrecy, the KDC should revoke all the keys. For the revocation, the KDC changes the old slave key to a new slave key and makes a new master key through the new proposed master key management algorithm.



(a)The Conceptual Diagram of MKE.(b)when  $d_1$  should be updated to  $d'_1$ ,the other private places  $\{d_1...d_r\}$  are still valid without changes

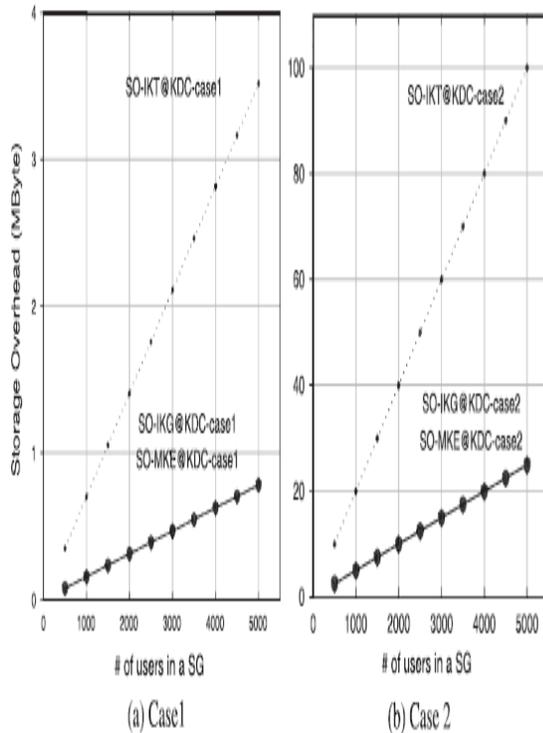


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SO at KDC for two cases



This case is applicable for hierarchical Service Set.

*The Proposed Master Key Management For Algorithm*

1. To determine  $p_1 \dots p_r, q_1 \dots q_r$  for safe prime no
2. For  $i=1 \dots r$
3.  $\alpha_i = (p_i - 1) * (q_i - 1)$
4.  $x_i = (p_i - 1) / 2$
5.  $y_i = (q_i - 1) / 2$
6.  $e_i = (4 * \text{random}) + 1$
7.  $d_i = e_i^2(x_i - 1)(y_i - 1) \text{ mod } (\alpha_i, y_i)$
8. end for
9.  $n = 1$

10. for  $i=1$  to  $r$
11.  $n = n * (x_i y_i)$
12. end for
13. for  $i=1$  to  $r$
14.  $M[i] = n / (x_i y_i)$
15.  $N[i] = M[i]^{(x_i - 1)(y_i - 1)} \text{ mod } (x_i y_i)$
16. end for
17.  $em = 0$
18. For  $i=0$  to  $r$
19.  $em = em + (e_i * m[i]^n[i]) \text{ mod } n$
20. end for
21. while  $(em \text{ mod } 4) \neq 1$   $em = em + n$
22. sleep
23. Interrupt (when  $j$ th key pair should be updated)
24.  $e_j = 4 * \text{random} + 1$
25.  $d_i = e_i^2(x_i - 1)(y_i - 1) \text{ mod } 4x_i y_i$
26. goto 17

## V. CONCLUSION

In this paper, MGKM has been proposed for multiple group instead of hierarchical group. In contrast to other existing system of using symmetric keys, the MKE-MGKM exploit asymmetric keys, master key and multiple slave key, which are generated from proposed master key generation. by using a set compromising master key and slave key, a TEK can efficiently distribute to multiple SGs. Therefore, number of rekeying efficiently reduced. Therefore keygraph of MKE-MGKM is much simpler than any other scheme, less memory is needed for storing the keys. Compared with other schemes, the MKE-MGKM can significantly reduce storage and Cos in the rekeying process, with acceptable computation overhead. It is expected MKE-MGKM is practical solution for various group application, especially those many SGs, especially TV charged services by channel basis.



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